

# A Communication Non-intrusive Middleware for Resource Management in Sparse MANETs

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# Overview

- Introduction
- Communication Non-intrusive Resource Management
  - A non-intrusive information source
  - Prediction
  - Clustering
  - Feasibility study
- Conclusions and Future Work

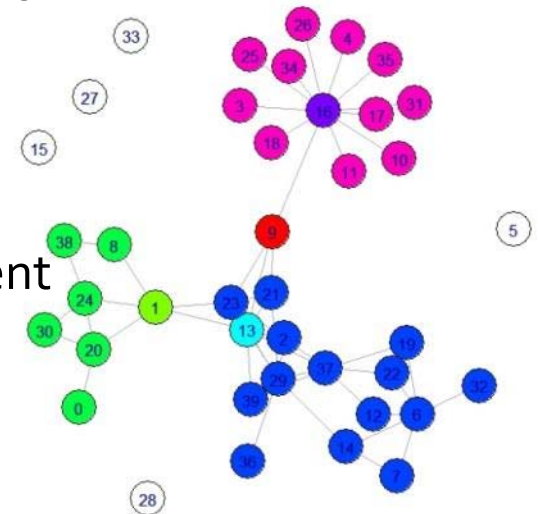
# Rescue Operations and Emergency Intervention



- Red Cross, 2005: “ ... the flow of information throughout the disaster cycle is crucial for effective humanitarian operations ... ”
- Resource Operation Specific:
  - Diverse scenes: landscape, area size, number of people, status of infrastructure, and time span
  - Heterogeneity: organizations, devices, and configurations
  - Cooperative personnel hierarchically organized
- Middleware for Rescue Operations and Emergency Intervention → Ad Hoc InfoWare middleware

# Ad Hoc InfoWare Middleware for Rescue Operations

- **Goal:** Improve the information sharing in rescue operations and emergency interventions
  - Increase the efficiency of collaborative work
- **Solution:** use data networks
  - MANETs + DTNs  $\approx$  Sparse MANETs
  - Advantages:
    - Infrastructure independent and fast deployment
    - Mobile and heterogeneous devices
  - Disadvantages
    - Mobile and heterogeneous devices
    - Unpredictable scene layout, environment, and movements of users  
→ unpredictable connectivity

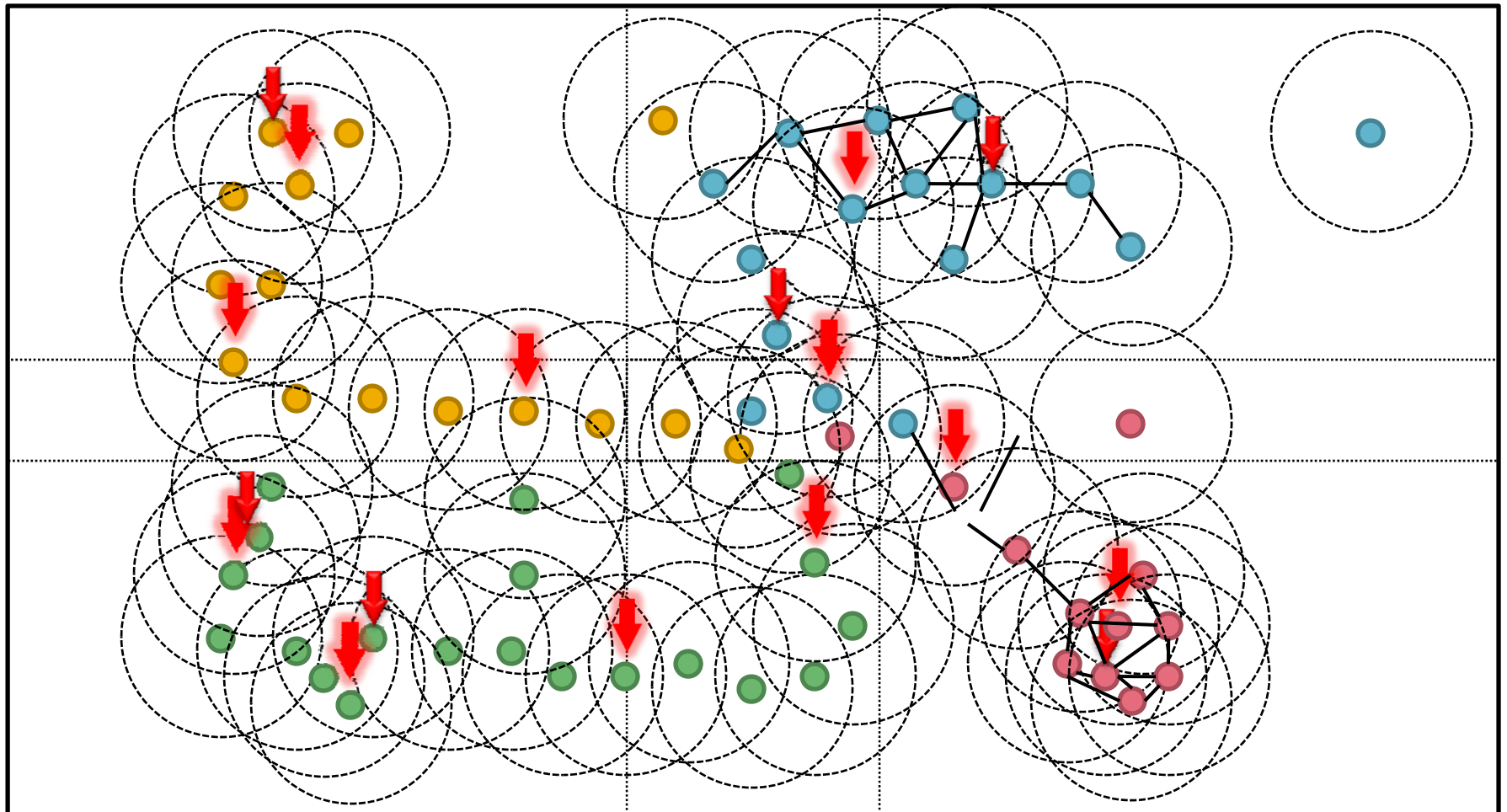


# Ad Hoc InfoWare Middleware – Resource Management

- Provide a network resource aware view → well-connected regions and possible partitions of the network
- *Resource Management*
  - *Balance* the use and allocation of resources
  - *Eliminate* communication overhead and dependence on external services
  - Provide applications with information on remote resource *availability*

Enable adaptation to changing network configuration and capabilities of devices.

# Resources & Service Placement – Example



# Claims

1. *Non-intrusiveness*
  - Routing protocol holds updated information
2. Neighborhood *prediction*
  - Past adjacency information
  - No constraints
3. Network *clustering*
  - Current network layout
  - No constraints
4. *Feasible* on resource constrained devices

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# Requirements analysis for Resource Management

- General assumptions
  - Set-up phase for devices and personnel
  - Heterogeneous devices → different capabilities and information
- Requirements
  - General solution (independent of routing protocol, device resource, and external service)
  - Local information
  - Preserve resources
  - Update resource information
- *Worst case scenario: no external and additional information*

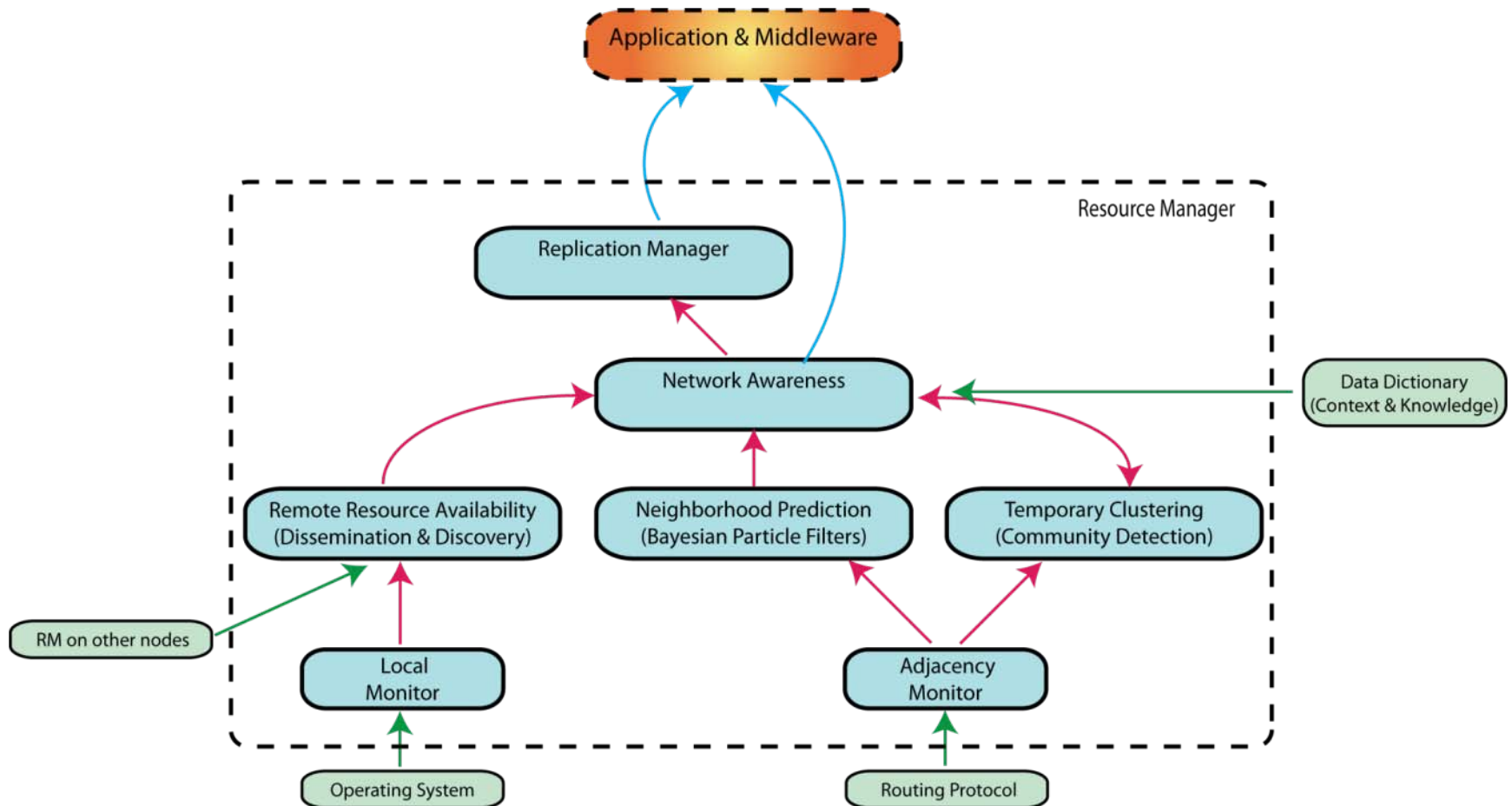
# Issues for Resource Management

- **Goal:** Enable composability:
  - Remote resource sharing (*Horizontal*)
  - Local resources (Vertical )
  - Resource sharing profiles (Hierarchical)
- Remote resource sharing
  - Reliable communication channel → Discovering, Bookkeeping, Negotiation
  - Un-reliable communication channel → Resource Availability (e.g., link-lifetime)
    - Graceful degradation, Resource planning, Timed-soft-state-like reservations

# Resource Management

- **Algorithm: Link-lifetime Estimation**
  - Classical: Global Positioning System (GPS)
    - Problems: imprecise measurements or lack of GPS coverage, physical proximity is not a guarantee of connectivity, needs information dissemination
  - Past Link-lifetime
    - Advantages: local information with no assumption on link-lifetime duration
- **Algorithm: Network Clustering**
  - Classical: Predefined number of clusters or cluster size
    - Problems: Dynamic of the network (number of nodes, connectivity, etc.)
  - Community Detection
    - Advantages: local information with no assumption on number of nodes in a cluster, number of clusters, network size and layout
- **Pros:** no assumption on location, no communication overhead
- **Cons:** more complicated algorithms, dependent on the accuracy of the information source

# Information flow of Resource Management



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# Foundation for a non-intrusive RM

- Information Source: Routing table in the routing protocol
  - Advantages
    - Updated view of the network,
    - Location independent
  - Disadvantages
    - Sensitive to existence of communication
    - Sensitive to mobility and communication patterns
    - Partial topology of the network
  - Information type:
    - Neighborhood
    - Topology
  - Issues: accuracy and consistency

# Test environment

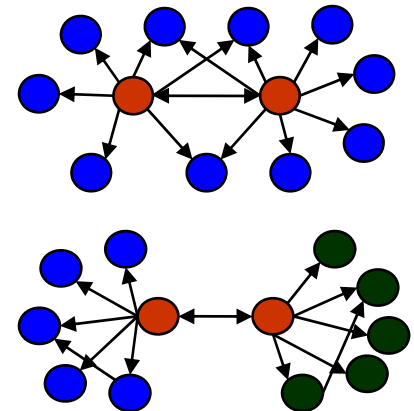
<u>Mobility Model</u>	<i>Single mobility models:</i> Random Direction Mobility, Random Walk Mobility, Steady-State Random Waypoint <i>Group mobility models:</i> Random Point Group Mobility <i>Other simulators:</i> <b><u>RoboCupRescue</u></b> Simulation Project	
<u>Network simulator</u>	<i>GloMoSim</i>	Reactive Routing protocol: <b>AODV</b>
		Proactive Routing Protocol: <b>OLSR</b>
<u>Communication</u>	<i>Flat:</i> a random subset of the nodes are transmission sources <i>Structured:</i> follows the group structure	

## General

- 40 nodes (4 groups)
- A=800x600u, T=18000s
- radius=100u,
- item=2048b,
- v=1÷7u/s

## Communication

- Flat: tr\_src>25%, per=50s
- Structured:
  - Leaders: tr\_src=10%, per={30s,60s}
  - Group Nodes: tr\_src>50%, per=30s



# Neighborhood Connectivity Periods

- **Issue:** Accuracy of neighbor information in route table
- **Question:** Are neighbor periods similar in the scene (mobility model) and the routing protocol?
- **Solution:** Discrete time (routing protocol) approximate continuous time to time periods: Neighbor or Non-Neighbor  $(n_1, n_2)$

- Aggregate timestamps of direct contact between a pair of nodes

- Interval  $[0, T_{sim}]$  where  $0 \leq t_1 < \dots < t_k \leq T_{sim}$

- $T_{sup} = \{t_{i+1} | t_{i+1} - t_i > \sigma\} \cup \{T_{sim}\}$   $T_{inf} = \{0\} \cup \{t_i | t_{i+1} - t_i > \sigma\}$   $\sigma$  - time threshold

- Neighbor time periods

$$C = \left\{ [t_i, t_j] \mid t_i \in T_{inf}, i = \{1, \dim(T_{inf}) - 1\}, t_j \in T_{sup}, j = \{2, \dim(T_{sup})\} \right\}$$

- Non-neighbor time periods

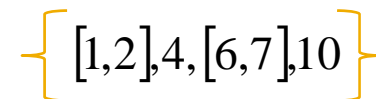
$$D = \left\{ (t_i, t_j) \mid t_i \in T_{inf}, t_j \in T_{sup}; i, j = \{1, \dim(T_{inf})\} \right\}$$



Timestamp  
Neighbor



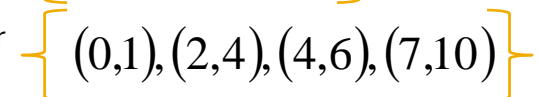
Neighbor  
Period



Timestamp  
Non-Neighbor



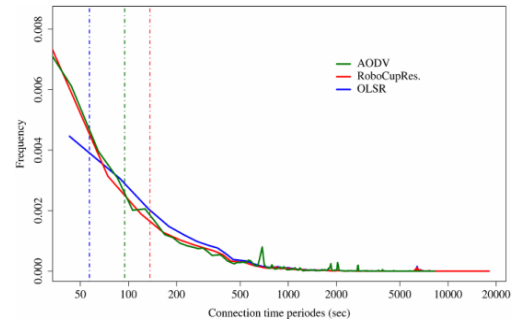
Non-Neighbor  
Period



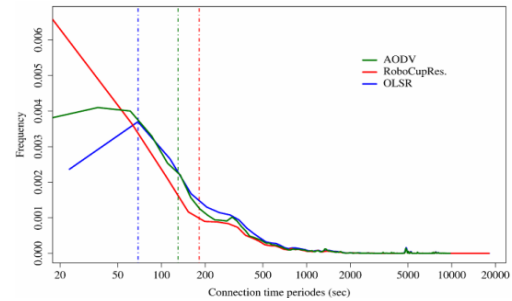


# RoboCupRescue – Distribution of Connectivity Periods

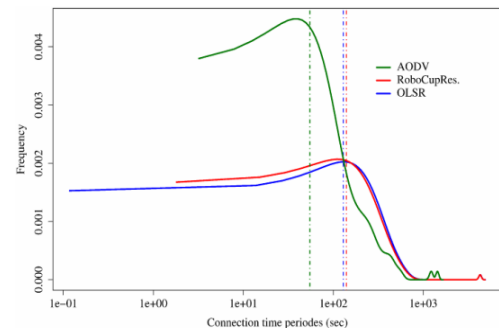
**Result:** The neighbor periods have the same distribution for mobility models and routing protocols.



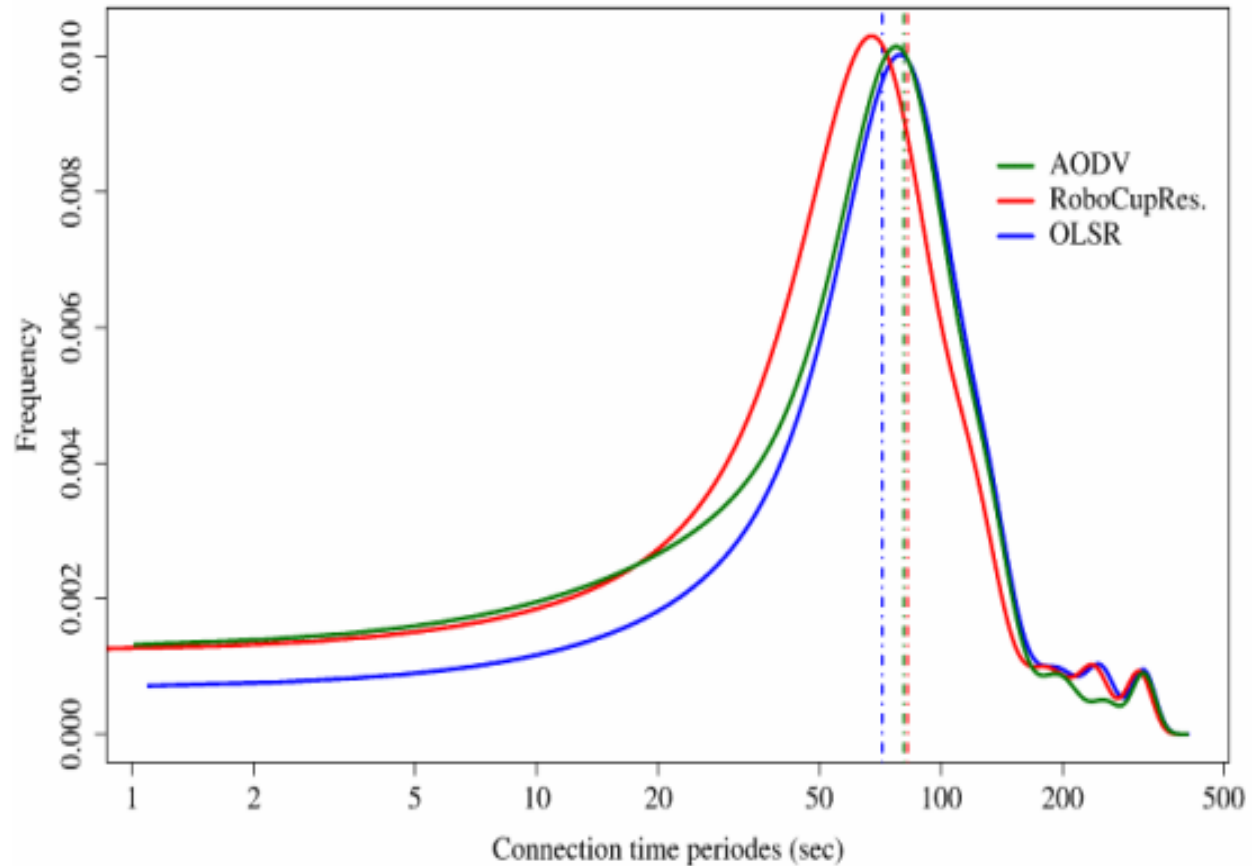
Connections – Flat



Connections – Struct.



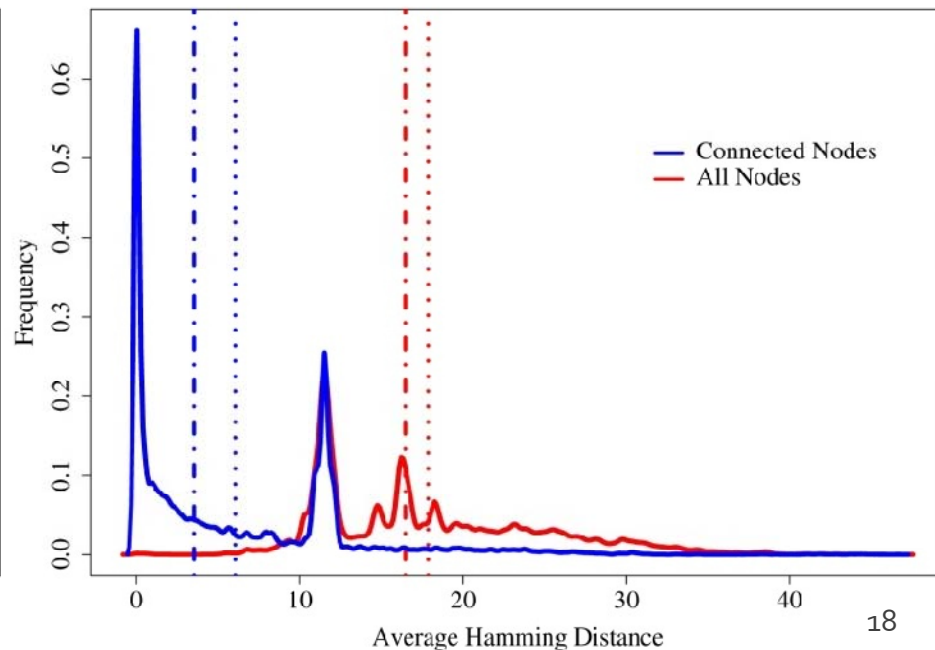
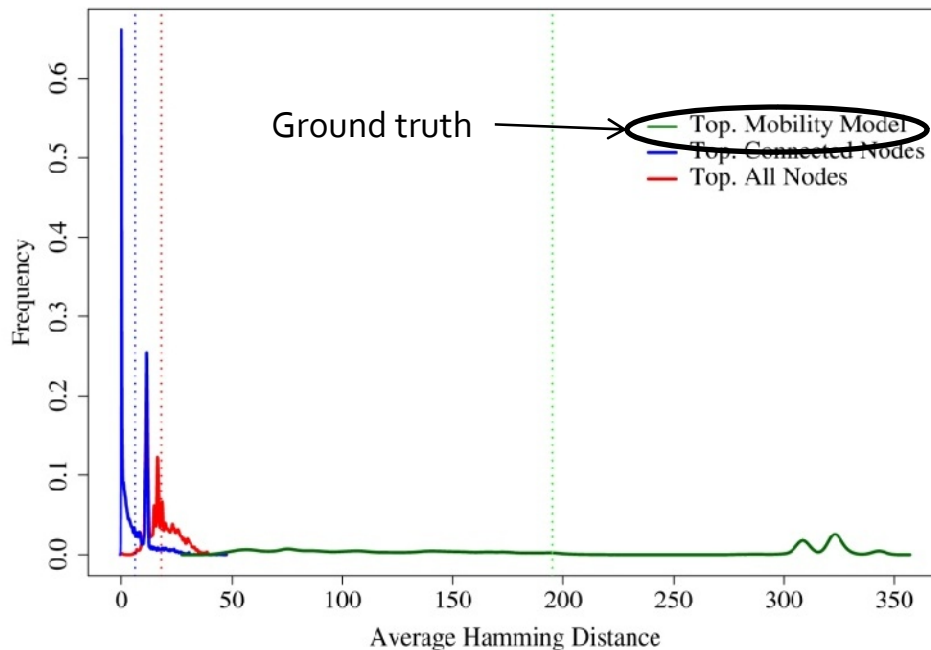
Group Connections – Flat



Group Connections – Structured communication

# Topology Data Consistency

- **Issue:** Topology consistency (i.e., nodes may have different topology information)
- **Question:** How consistent is the topology information at the different nodes?
- **Solution:** Compare topology information on all the node in the network
  - Compute the Hamming Distance between topologies (count the differences between the adjacency matrixes of different nodes)
- **Result:** Similar topologies, if node are connected.



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# Prediction Algorithm

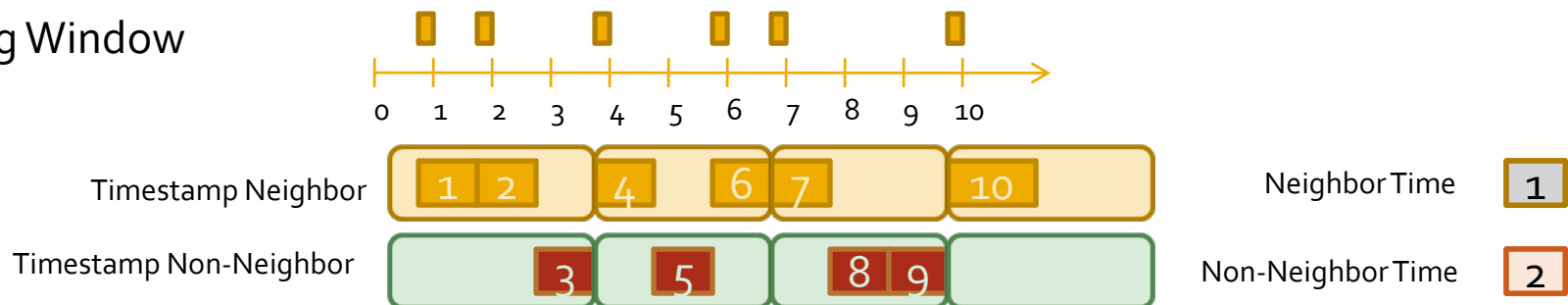
- **Question:** Is the neighborhood of a node predictable?
- **Idea:** Neighborhood estimation algorithm based on routing information
  - Active communication → routing protocol has good information on neighborhood
  - Discrete timestamps approximate continuous time → Neighbor or Non-Neighbor Periods
- **Algorithm:** SIR Future Neighbor
  - Window in a data stream (data process model)
    - Changes of direct connectivity between pairs of nodes
  - Sequential Monte Carlo (estimation method)
    - Estimates the future state of a node's neighborhood
    - No assumptions on: Links' lifetime and Nodes' movement

# Window in a data stream

- Solution: Neighborhood window**
  - Mapping from the time domain  $T$  to the domain of possible time intervals  $T_s$ 

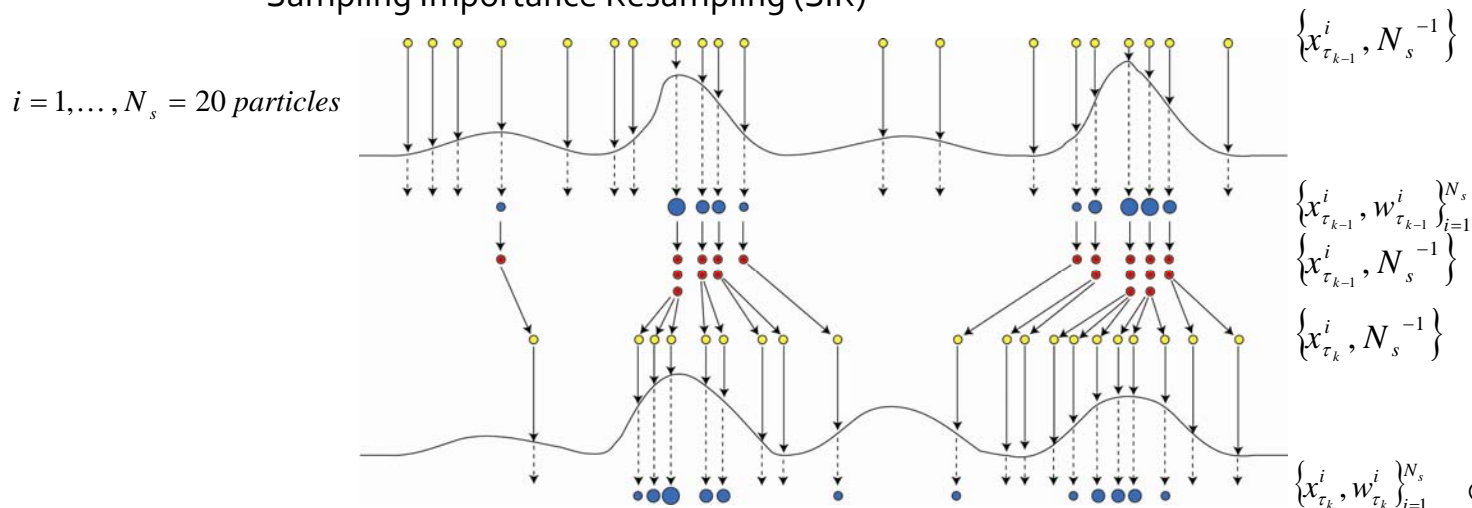
$$\Psi : T_s \rightarrow \{[t_1, t_2] \mid t_1, t_2 \in T, t_1 \leq t_2\}$$
  - Aggregate timestamps of direct contact between two nodes
    - State of the pair  $(n_1, n_2)$  during an interval  $[t_j, t_{j+1}]$  where
 
$$\min(\psi(t)) \leq t_1 < \dots < t_k \leq \max(\psi(t)) \text{ and } \psi(t) = [\min(\psi(t)), \max(\psi(t))]$$
  - Neighbor time period
 
$$C_{\Psi(t)}^{(n_1, n_2)} = \sum_{j=1}^{k-1} (t_{j+1} - t_j) \text{ if } t_{j+1} - t_j \leq \sigma \text{ } \sigma\text{-time threshold}$$
  - Non-neighbor time period
 
$$D_{\Psi(t)}^{(n_1, n_2)} = (\max(\Psi(t)) - \min(\Psi(t))) - C_{\Psi(t)}^{(n_1, n_2)}$$
- Assumptions:** Distributions of periods are non-linear / non-Gaussian

Tumbling Window

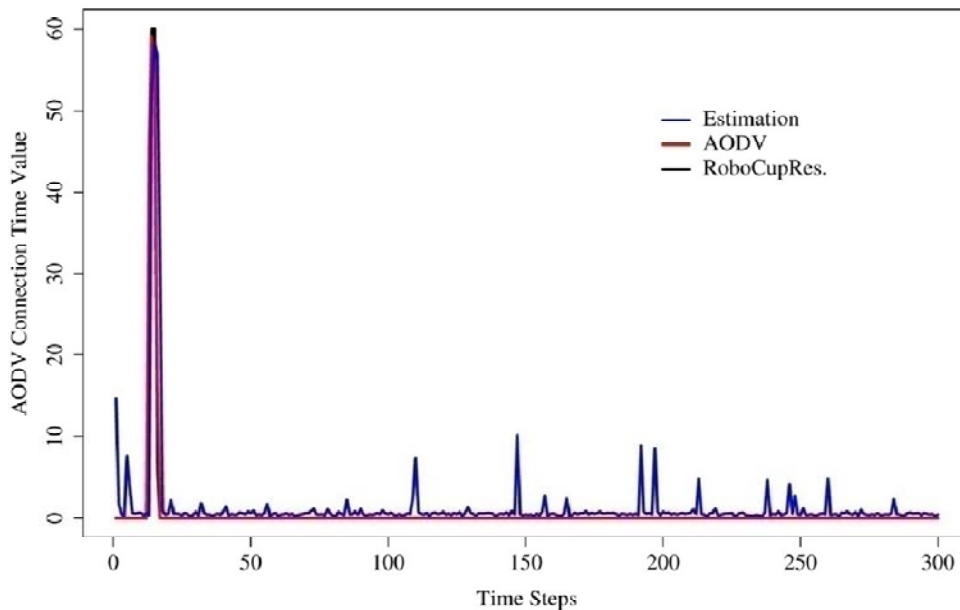


# Sequential Monte Carlo (SMC)

- *Recursive Bayesian Tracking*: Represent the posterior density function by a set of random samples with associated weights and compute the estimates based on samples and weights
- SMC for the neighbor pair  $(n_1, n_2)$ 
  - Recursively calculate the degree of belief in the estimation  $x_{\tau_k}$  for the next window interval  $\tau_k$  taking into consideration the values of the current measurements  $z_{\tau_{1:k}}$ 
    - Predict: use previous measurements for prior  $p(x_{\tau_k} | z_{\tau_{1:k-1}})$
    - Update: likelihood function  $p(z_{\tau_k} | x_{\tau_k}) \rightarrow$  prior distribution  $p(x_{\tau_k} | z_{\tau_{1:k}})$ 
      - Sampling Importance Resampling (SIR)



# Tumbling Window – SIR Future Neighbor



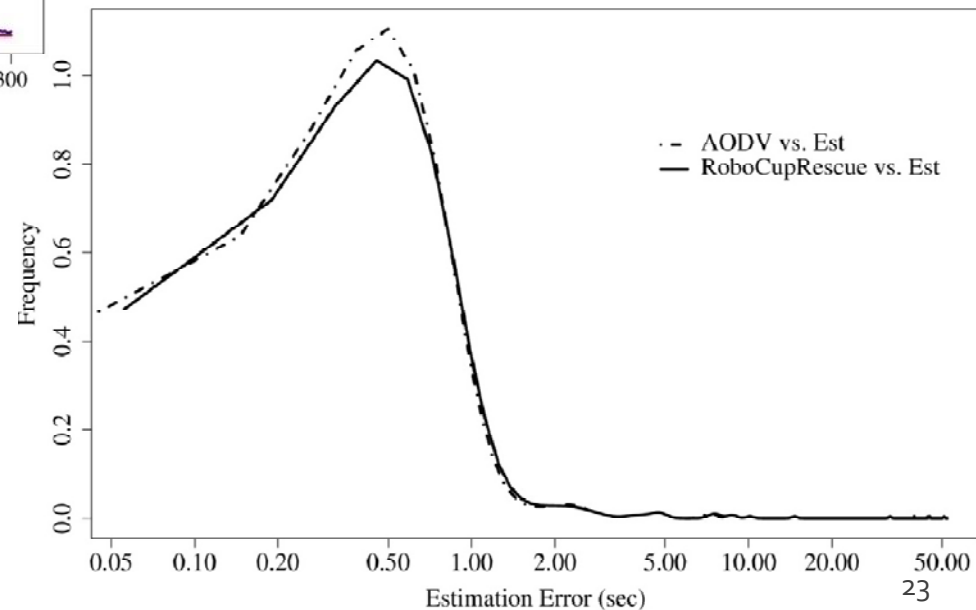
Connection time

Window type: Tumbling  $\rightarrow$  Estimate a value between 0 and  $\phi$

$$C_{\Psi(t)}^{(n_1, n_2)}, D_{\Psi(t)}^{(n_1, n_2)} \in [0, \phi]$$

$$\phi = 60$$

Euclidian Distance



**Result:** The estimated values are close to the measured and "true" values.

# Prediction Discussions

- Window Size ( $\phi$ ) vs. Time Step ( $\beta$ )
  - Gives the type of window, Impacts the performance and resources required,  $\beta$  does not influence accuracy
- Reactive routing protocol
  - No communication  $\rightarrow$  No accuracy
  - Windowing techniques can filter small communication pauses
- Linear Cost with Number of Samples ( $N_s$ )
  - Number of Samples ( $N_s$ ) and Time Step ( $\beta$ ): Large  $\beta \rightarrow$  Large  $N_s$
  - Number of Samples ( $N_s$ ) and Number of neighbors: High number of neighbors  $\rightarrow$  High computation cost
- Run continuously
  - Adapts to the behavior of the neighborhood, Correct after a few estimations
  - No on-demand estimations

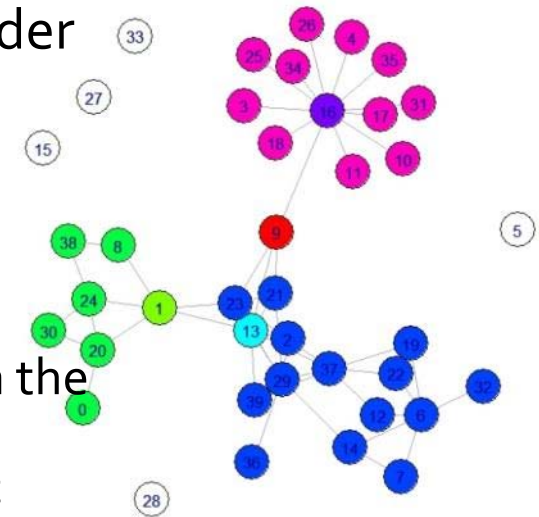


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# Temporary services

- **Question:** Where to place services in the network?
- **Issue:** Minimize the distance to resources in order to balance the use of resources
  - *Requirements*
    - Management overhead independence
    - Position independence
  - *Solution:* Temporary Clustering with dynamic clustering methods (i.e., consider the dynamic in the network)
    - Clustering which adapts to the current network layout
      - Adaptive number of clusters
      - Unconstrained number of nodes in a cluster
    - Temporary service positioning: Number of data replica and services, Data and service placement, Network partitioning
  - *Problem:* No methods to handle dynamics issues of MANETs



# Passive Clustering

- **Idea:** Clustering based on the network topology in the route table
- **Methods:** use static networks algorithms
  - *Clustering:* Community Detection
    - Divide or agglomerate to detect the groups of nodes in the network with dense network connections, and sparse connections outside the groups
    - 1. Newman and Grivan 2004 (Recursive ) where recursively eliminates the links with high weight in the network
    - 2. Reichard and Bornholdt 2006 (SpinGlass ) where community membership of a node is determined by its neighborhood (i.e., number of neighbors and neighbors' membership)
  - *Placement:* Cluster head election based on centrality measures
  - *Coverage:* Network Voronoi Node Diagram
    - For nodes in a given set it defines: dominance set and bisector set nodes
  - *Evaluation:* measure quality, stability, similarity, and consistency

# Clustering Evaluation: Quality

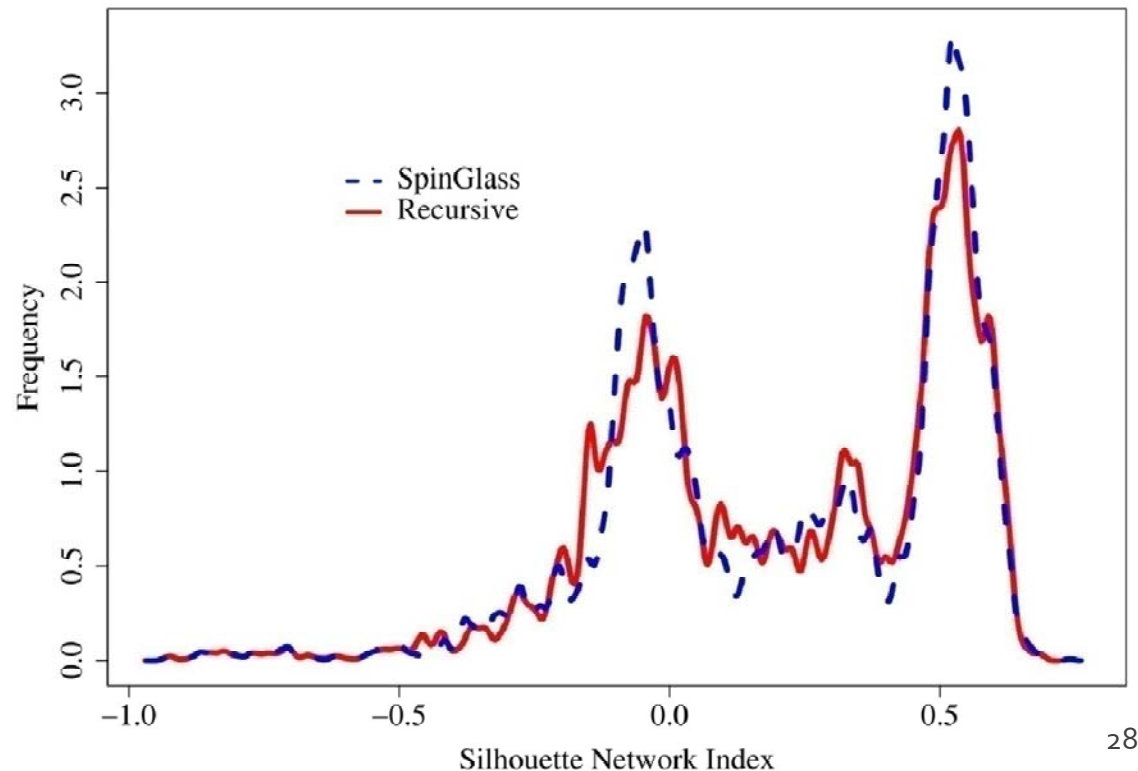
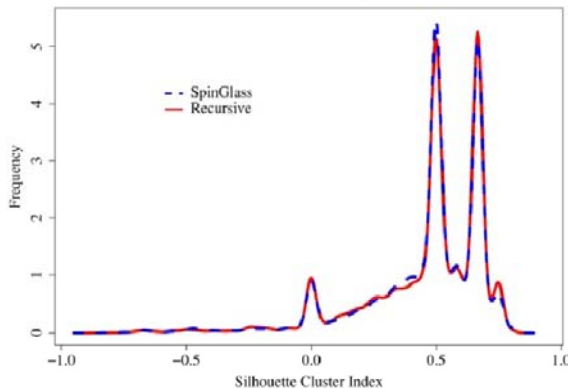
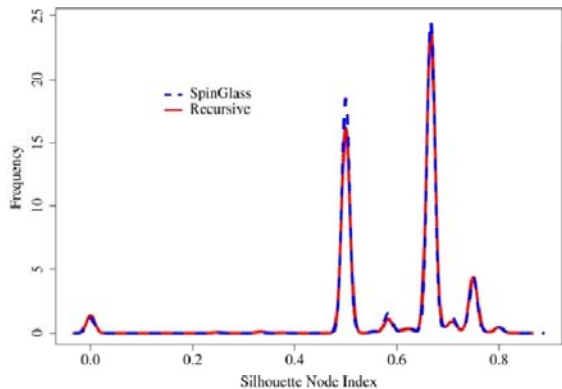
- **Question:** Is the clustering valid?
- **Measure:** Silhouette index (how well is a node clustered considering its distance to the center and of the center of the closest cluster)

Node:  $s(p_i) = \frac{\bar{d}(p_i, C_h) - \bar{d}(p_i, C_j)}{\max(\bar{d}(p_i, C_h), \bar{d}(p_i, C_j))}$

Cluster:  $S_j = \frac{\sum_{i=1}^{N_j} s(p_i)}{N_j}$

Network:  $GS = \frac{\sum_{j=1}^k S_j}{k}$

- **Results:** The created clusters are good.

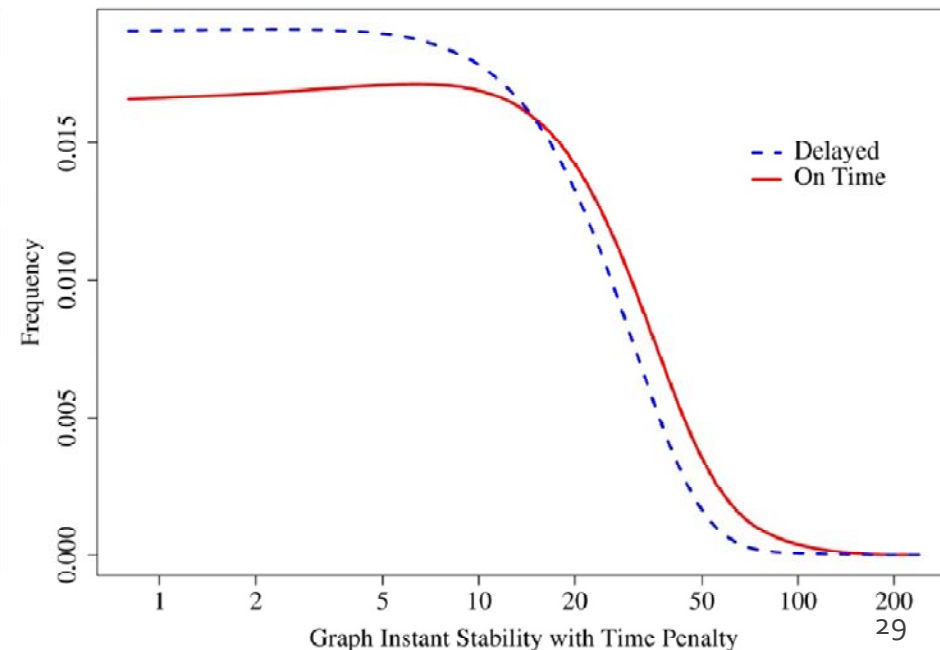
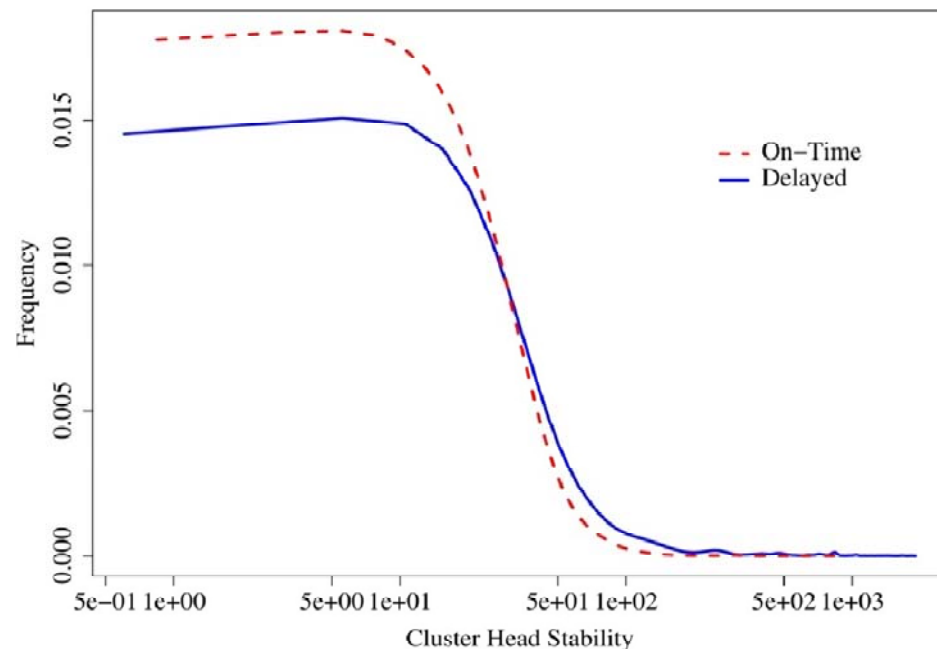


# Clustering Evaluation: Stability

- **Question:** Is the clustering stable?
- **Measure:** Stability quantifies the changes of the clustering with respect to the new network structure (i.e., establishment and breakage of links, and quality increase and decrease cluster membership)

$$S_i = \left[ \frac{Q_i + F_i}{B_i + E_i} \right] * \frac{1}{\delta t}$$

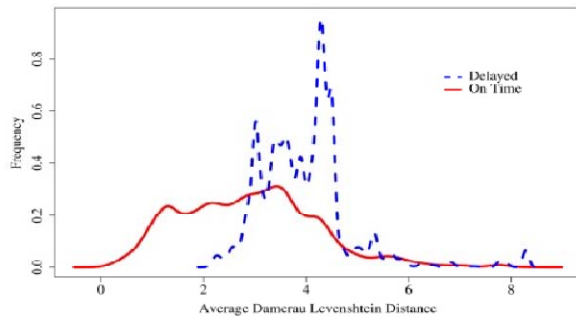
- Delay the clustering
- **Results:** Delayed clustering can improve the stability of the clusters.



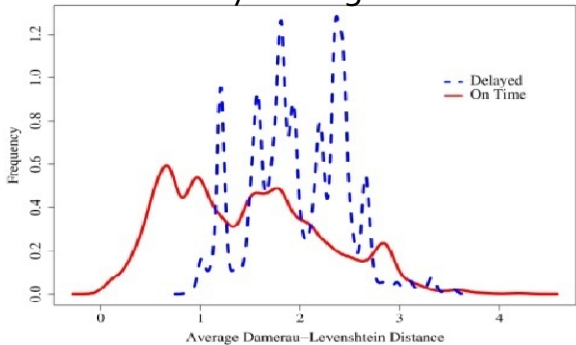
# Clustering Evaluation: Similarity

- **Question:** What is the difference between different clusterings?
- **Measure:** The similarity measures the variation of information between clustering over the same network
  - SpinGlass vs. Recursive
    - The methods return similar clustering, i.e., in the interval 0 and 0.8
  - SpinGlass vs. Voronoi and Recursive vs. Voronoi
    - The Voronoi Diagrams are also similar to the clustering's, i.e., in the interval 0.1 and 0.9
- **Result:** The different clustering techniques produce similar results.

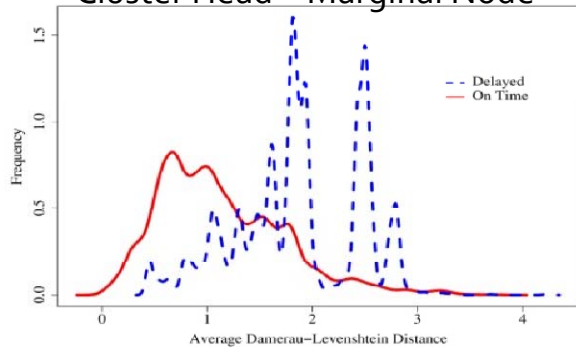
# Clustering Evaluation: Consistency



Community – Marginal Node

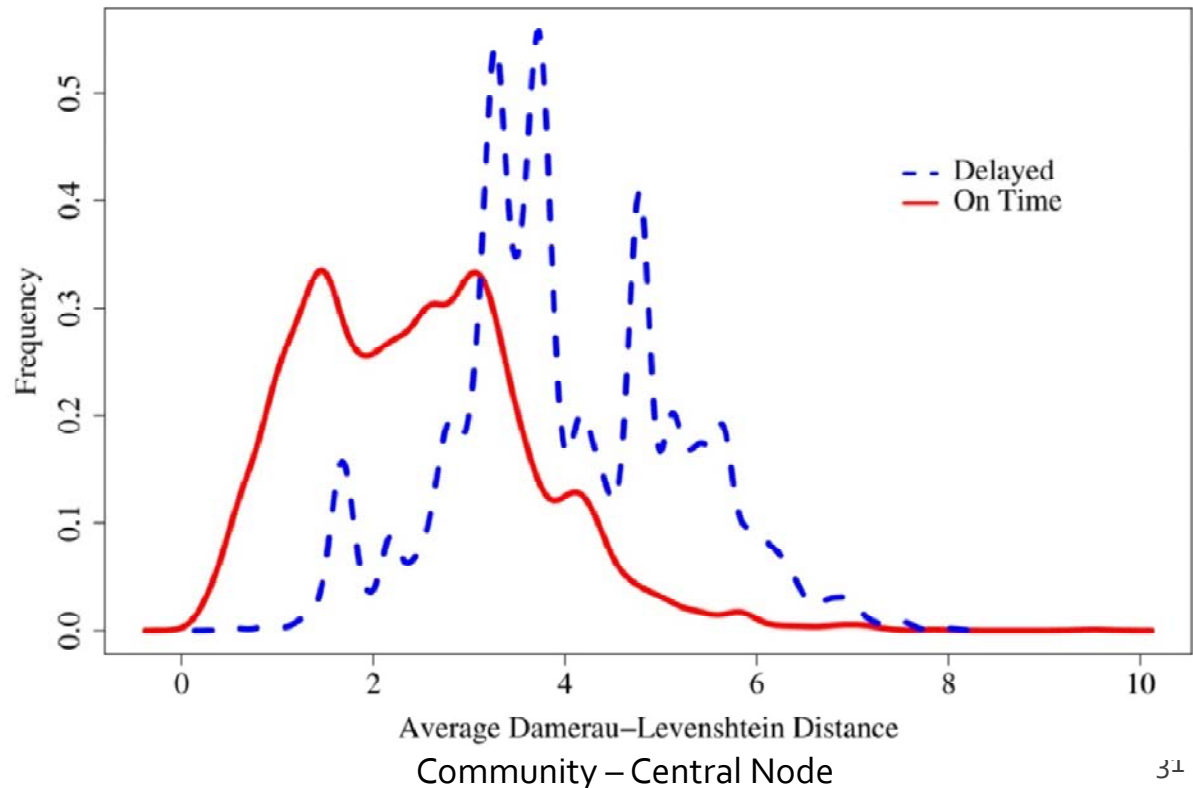


Cluster Head – Marginal Node



Cluster Head – Central Node

- **Question:** Are the clusters consistent in the network?
- **Measure:** Damerau-Levenshtein Distance between detected communities at different nodes (i.e., no. of insertions, deletions, substitutions of single characters, and transpositions between two sets)
- **Results:** Communities are similar at different nodes.



# Clustering Discussion

- Topology information
  - Does not apply for reactive routing protocols
  - Requires a consistent view of the topologies at the nodes
- Clustering measures
  - Not for dynamic networks
  - Quality
    - Not a general accepted metrics
    - Different metrics may give contradicting conclusions
  - Stability
    - Does not consider the changes in the number of clusters
    - Does not consider the changes in the number of nodes in the network
  - Similarity:
    - Not applicable to clustering from different nodes in the network
  - Consistency
    - Detected communities are more consistent than elected cluster heads



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# Implementation – Nokia 770 Internet Tablet

- SIR Future Neighbor
  - Values
    - Footprint 28kb
    - Routing protocol: OLSRD
    - $\beta = \phi = \{30, 60, 120, 300\}$
    - $N_s = 20, T = 5200$
    - Number of nodes = 5 – 25 ,  
Simulated neighborhood  
change every 4 s
  - Computation time per step
    - Tumbling Window process ~  
0.03 – 2 ms
    - SIR Future Neighbor process  
~ 2 – 40 ms

**Result:** Possible to run on resource constraint devices.



*Image from [www.nokia.com](http://www.nokia.com)*

- Network Voronoi Node Diagrams
  - Values
    - Footprint 38.5kb
    - Routing protocol: OLSRD
    - Number of nodes = 10 – 15 ,  
Number of clusters = 2 – 4

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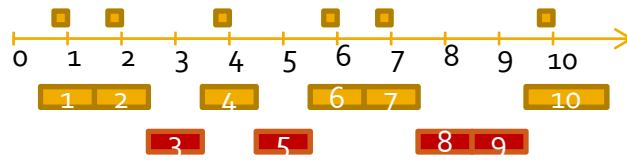
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# Conclusions

- Thesis contributions
  - Middleware architecture for RM in Sparse MANETs
  - Information service for worst case scenario to deliver updated availability of remote services and resources
  - Communication non-intrusive algorithms for prediction and clustering
- Thesis claims: Critical review
  1. Non-intrusiveness
    - Good if the node is involved in communication
  2. Neighborhood prediction
    - Non-intrusive is possible
    - Dependent on communication
    - Cannot perform timely predictions
  3. Network Clustering
    - Community detection gives good results
    - Can benefit from delaying the process
    - Only for proactive protocols
  4. Feasible on resource constrained devices

# Future Work

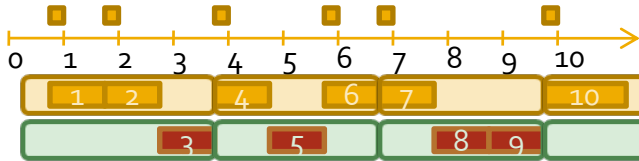
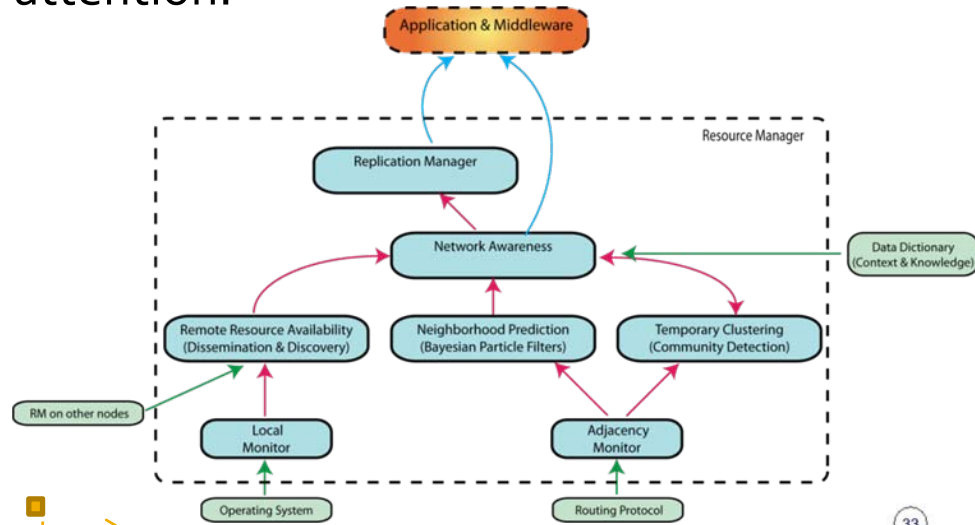
- Integrate higher level information
- Timely predictions
- Clustering for reactive protocols
- Use predictions and adaptive clustering techniques (which support the dynamics of the network)
- Apply the algorithms to delay tolerant streaming
- Evaluate with real mobility and communication traces
- Evaluate during a rescue exercise



Neighbor Period  $\{ [1, 2]_4, [6, 7]_{10} \}$   
 Non-Neighbor Period  $\{ (0, 1), (2, 4), (4, 6), (7, 10) \}$

Thank you, for your attention!

Questions?



Neighbor Time  $\{ 1 \}$   
 Non-Neighbor Time  $\{ 2 \}$

